

FIG. 1

	162
/	

Tag	Operation <u>166A</u>	Operation <u>166B</u>	Operation <u>166C</u>	Operation <u>166D</u>	Flow Control 168
<u>164</u>	Operation <u>166E</u>	Operation <u>166F</u>	Operation <u>166G</u>	Operation <u>166H</u>	

FIG. 2A

Liveness	Encoding 266
Unconditionally Live	11
Subsequent to 1st Branch	10
Subsequent to 2nd Branch	01
Unconditionally Dead	00

FIG. 2B

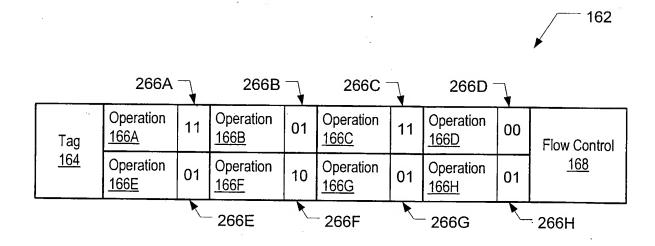


FIG. 2C

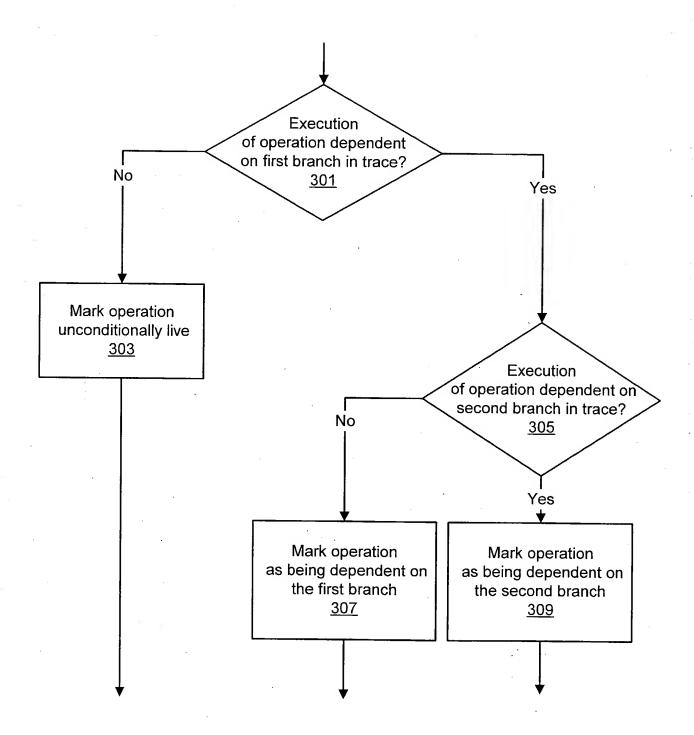
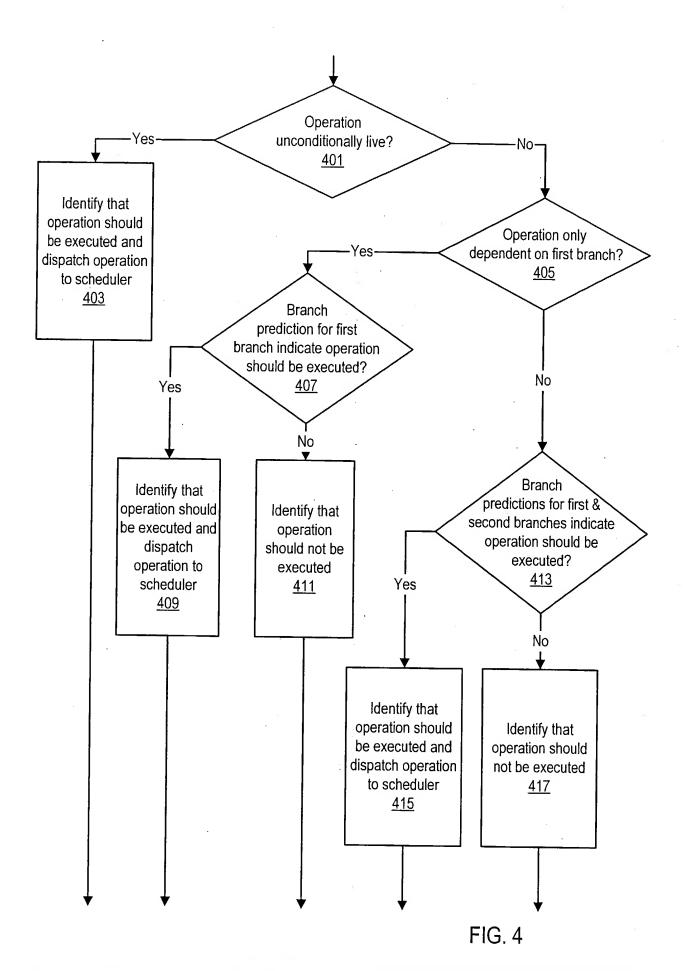
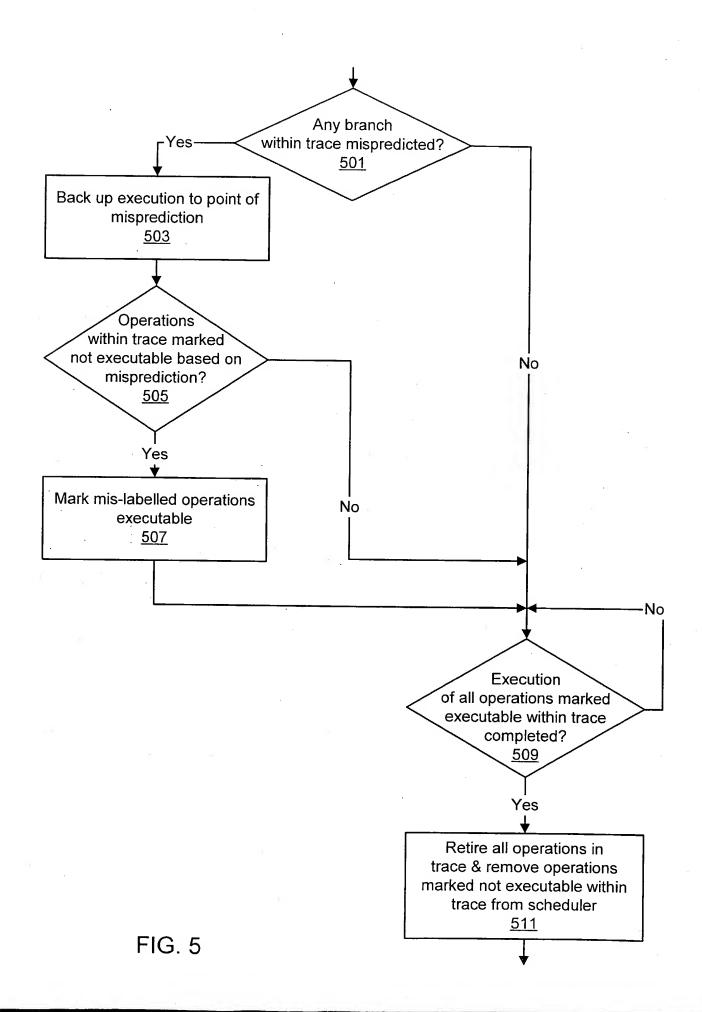
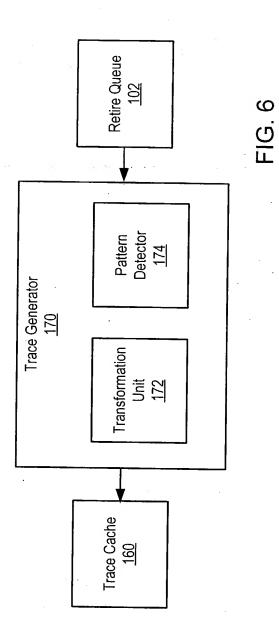


FIG. 3







Before Constant Propagation:

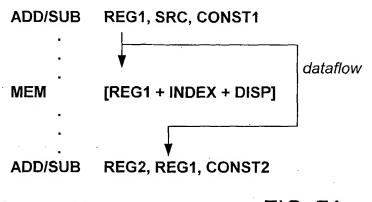
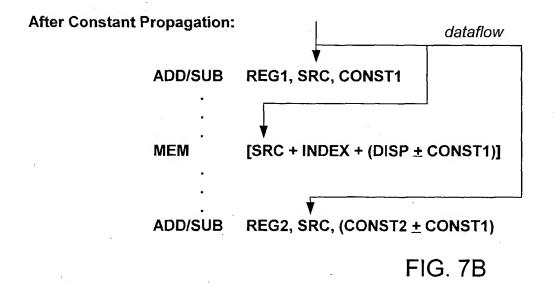


FIG. 7A



Before Move Renaming:

After Move Renaming:

MOV EAX, ESI

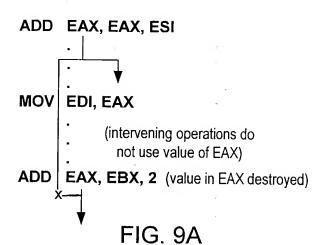
ADD EAX, ESI, 7

ADD EAX, EAX, 7

FIG. 8A

FIG. 8B

Before Move Renaming:



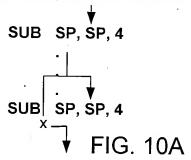
After Move Renaming:

ADD EDI, EAX, ESI

ADD EAX, EBX, 2

FIG. 9B

Before any transformations:



Slotification:

After Constant Propagation:

After Dead Operation Elimination:

FIG. 10D

Before:

Op1 Sets condition code

(intervening instructions do not set condition code)

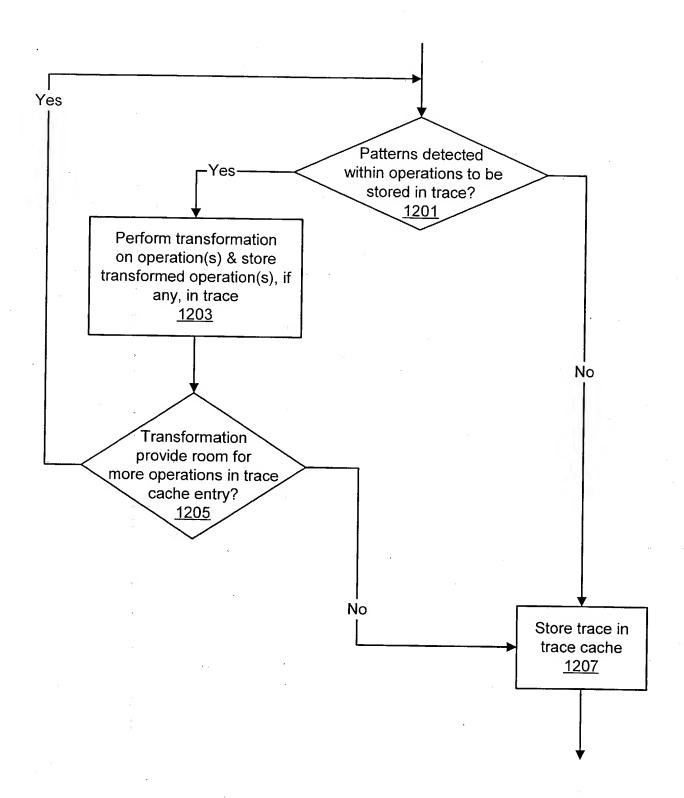
JNE Label

FIG. 11A

After Branch Folding:

Op1 Sets condition code; fail if NE

FIG. 11B



Instruction Stream:

Instruction1
Instruction2
Instruction3
Microcoded Instruction
Instruction4
Instruction5

FIG. 13A

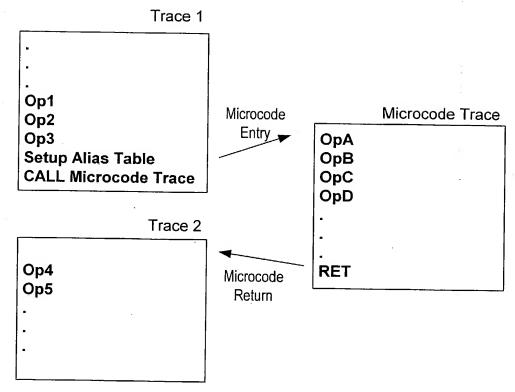


FIG. 13B

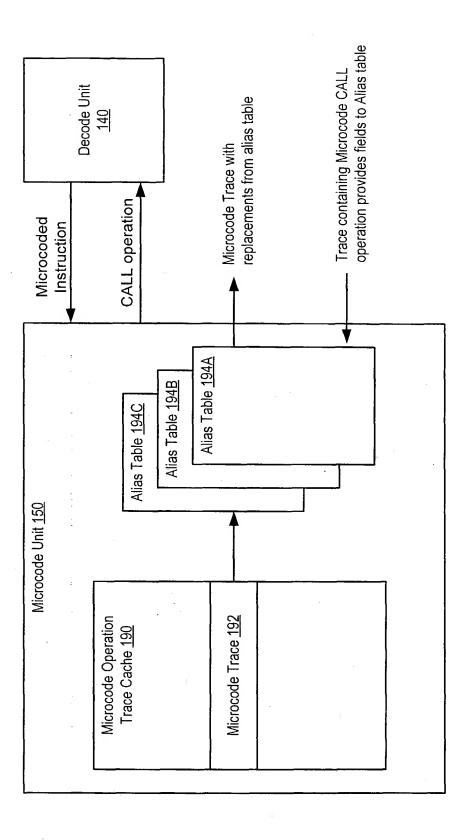


FIG. 13C

Replace microcoded instruction with an alias table setup operation and a call operation that calls the appropriate microcode subroutine and that specifies the appropriate alias table register names <u> 1401</u> Execute call operation, which pushes return address (identifying next trace to execute after execution of microcode subroutine) onto stack <u>1405</u> Read microcode subroutine operations specified by call operation and perform replacements using alias table register names as microcode subroutine operations are dispatched to schedulers <u>1403</u> Execute return operation included in microcode subroutine, which pops return address from stack 1407

